

Stack operation optimization for Michelson

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4/18 @ lab8 lunch talk

Background

- Michelson
 - for Tezos smart-contract development
 - Stack based
 - c.f.) Forth / Java bytecode / OCaml bytecode
 - Statically typed
 - Hard to write by hand
 - Various high-level languages & compilers are developed

Michelson code example

- <https://smartpy.io/ide>
 - Online IDE for Python → Michelson compiler

Michelson in Real World

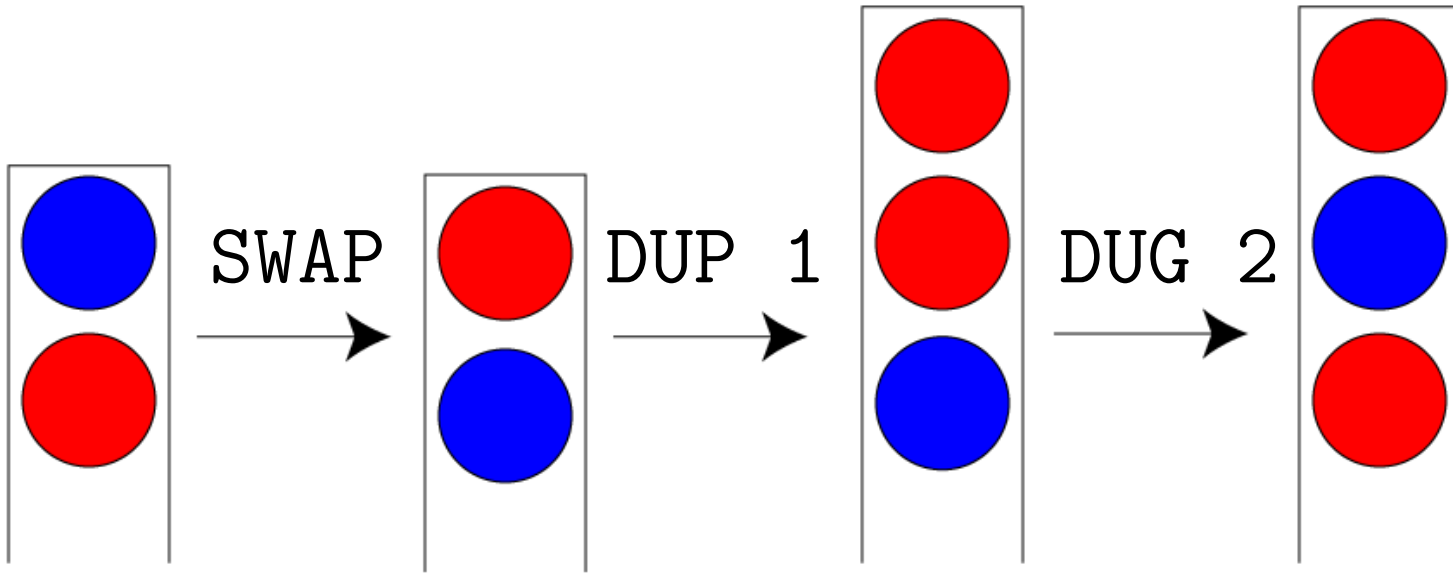
- Program costs "storage burn" in proportion to its size
 - 1byte ~ 0.001\$ ~ 0.1 yen
- Unoptimized contracts are deployed
 - Compiler emits unoptimized code
 - Optz (our optimizer) reduces avg. 5% of size

Optz

- <https://dailambda.jp/optz/>
 - Online editor: <https://dailambda.jp/optz-js/>
- Optz optimizes Michelson code in 3 ways
 1. pattern matching
 - `{ SWAP; LT; } → { GT; }`
 - `{ DROP n; DROP m; } → { DROP (n+m); }`
 2. Exhaustive search on stack manip op seq
 3. Special case optimization

Optimization Example of Exhaustive search

- Before) SWAP; DUP 1; DUG 2



- After) DUP 2

Target of exhaustive search

Sequence of stack manipulation operations

- Target
 - PUSH, DUP, DROP ... insert/delete elements
 - SWAP, DIG, DUG ... rearrange elements
- Non-target
 - ADD, MUL, CMP ... calculation
 - IF, LOOP ... control operators

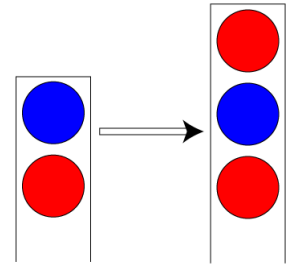
Stack manipulation function

- Stack manipulation operation sequence

represents stack \rightarrow stack function

- Symbolic execution result of stack

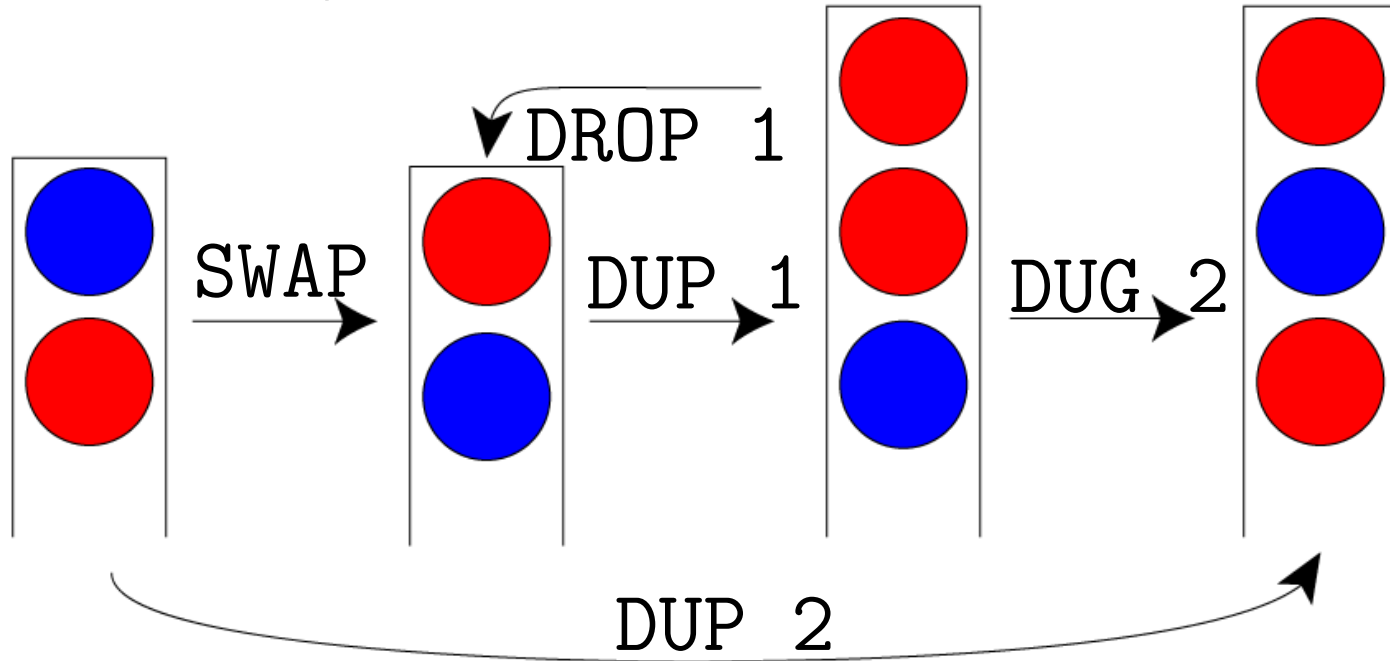
- e.g.) { SWAP; DUP 1; DUG 2; } ...
 { DUP 2; }



- Find cheapest stack op seq
which represents same function

Exhaustive Optimization using Dijkstra search

- Find best op with Dijkstra search
 - Cost of edge is `op size + op exec cost`



Computational Complexity of exhaustive search

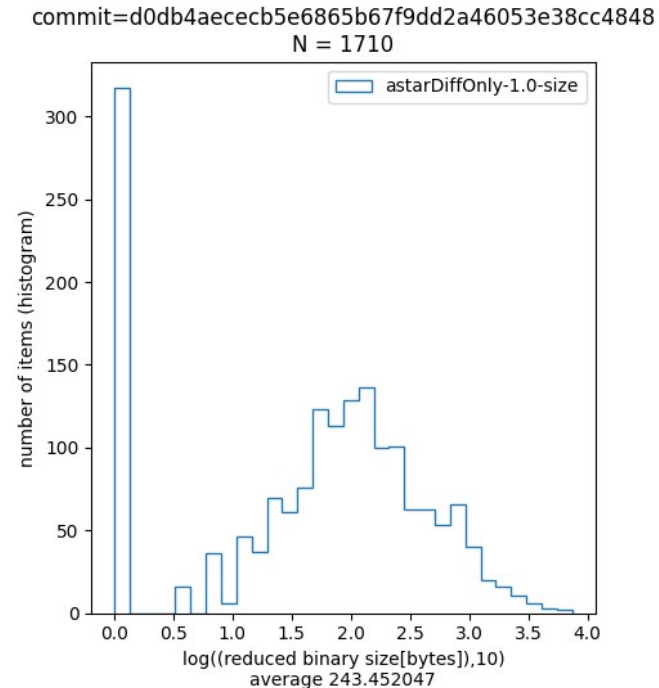
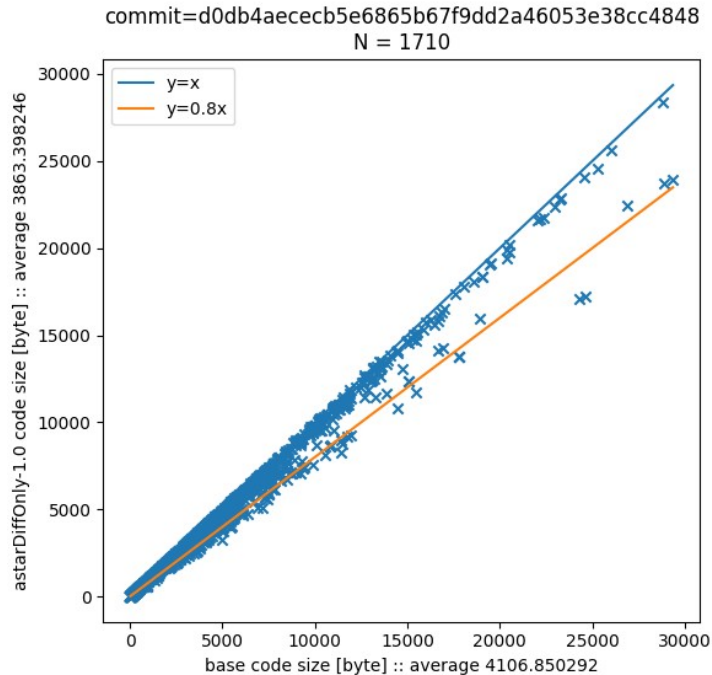
- $O(N^2 \log N)$ where N is "the num of nodes in graph"
- N is upper bounded by
 $L := \text{num_of_variable_variations} \wedge \text{max_stack_length}$
- We use L to prevent time-consuming optimization.
- Empirically, we set L to 10000.0
 - Every sampled contracts is optimized within 1[s].

Special Case Optimization

- Drop-only op seq
 - Result stack is sub-sequence of start stack
 - e.g.) { SWAP; DROP; SWAP; DROP; ... }
 - L is too big to optimize with exhaustive search. Instead, we use ad-hoc optimization.
 - Such unoptimized drop-only seq is generated from stack cleaning in function epilogue.

Result of optimized code size

Contracts deployed in 2021-02-18 - 2022-02-17
Avg. size reduction is 5% / 243byte (~ 24 yen)



A* search for faster optimization

- Dijkstra search
 - Search from the node whose `cost(node)` is the smallest.
- A* search
 - Estimate `score(node)` which satisfies
$$0 \leq \text{score}(\text{node}) \leq \text{actual_distance}(\text{node}, \text{goal})$$
 - e.g.) solving maze \rightarrow `score(node) := Manhattan distance to goal`
 - Use `cost(node) + score(node)` instead of `cost(node)` for searching

Good score function for stack modification op seq

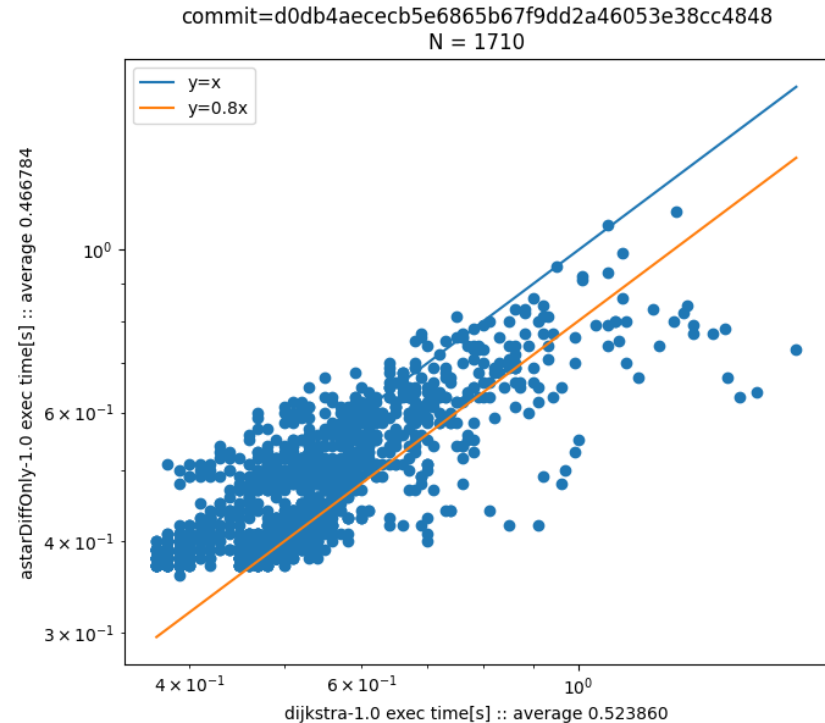
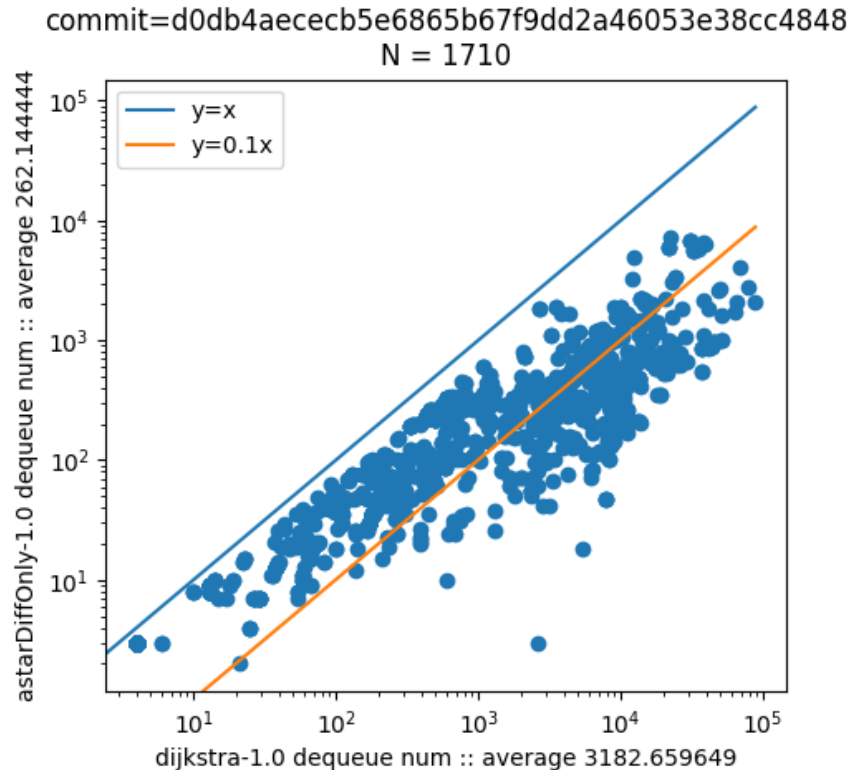
- Required number of PUSH operations
 - $\text{score}(\text{stack}) := \text{size}(\text{set}(\text{goal}) - \text{set}(\text{stack})) * \text{cost}(\text{PUSH})$
 - The vars appear in goal,
and not appear in current stack,
should be pushed.

Result of Optimization Speed

A* v.s. Dijkstra

Reduced Searched Node :: 90%

Reduced Time :: 10%



Summary

- Optz uses 3 types of optimizations, pattern matching, exhaustive search, and drop-only sequence.
- Optz reduces avg. 5% size of code.
- A* reduces 10% of exhaustive search time.

Appendix: Materials / References

- <https://dailambda.jp/optz/>
 - Blog post for optz
- <https://dailambda.jp/optz-js/>
 - Web interface for optz
- <https://gitlab.com/dailambda/scaml/-/blob/master/src/michelson/optimize.ml>
 - Source code of the optimizer part of optz
- <https://medium.com/hackernoon/optimizing-stack-manipulation-in-michelson-31ba7ff11a3a>
 - Source of the idea for using *A

Appendix: score function with LCS

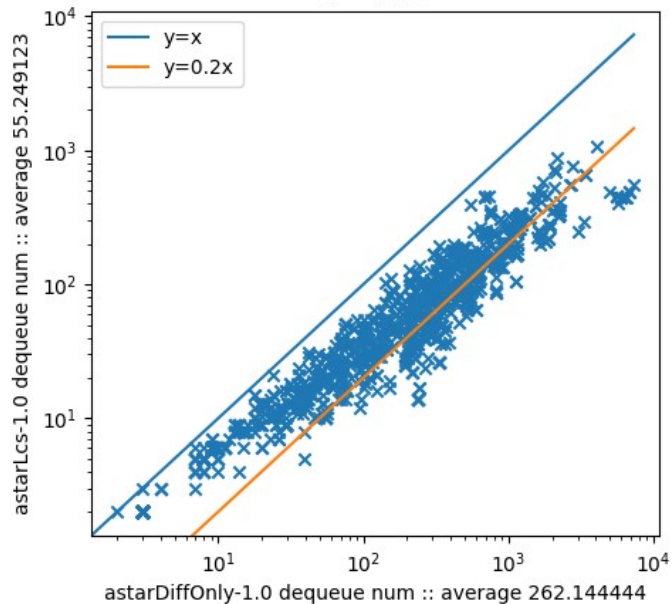
- Estimate with LCS
 - Good op sequence will preserve stack ordering
 - $\text{score}'(\text{st}) :=$
 - let $l = \text{LCS}(\text{st}, \text{goal})$ in
 - let $\text{push} = |\text{set}(\text{goal}) - \text{set}(\text{st})|$ in
 - $(\text{len}(\text{st}) - 1) * \text{cost}(\text{DROP}) +$
 - $(\text{len}(\text{goal}) - \text{push} - 1) * \text{cost}(\text{DUP}) +$
 - $\text{push} * \text{cost}(\text{PUSH})$

Appendix: LCS score v.s. PUSH diff only score

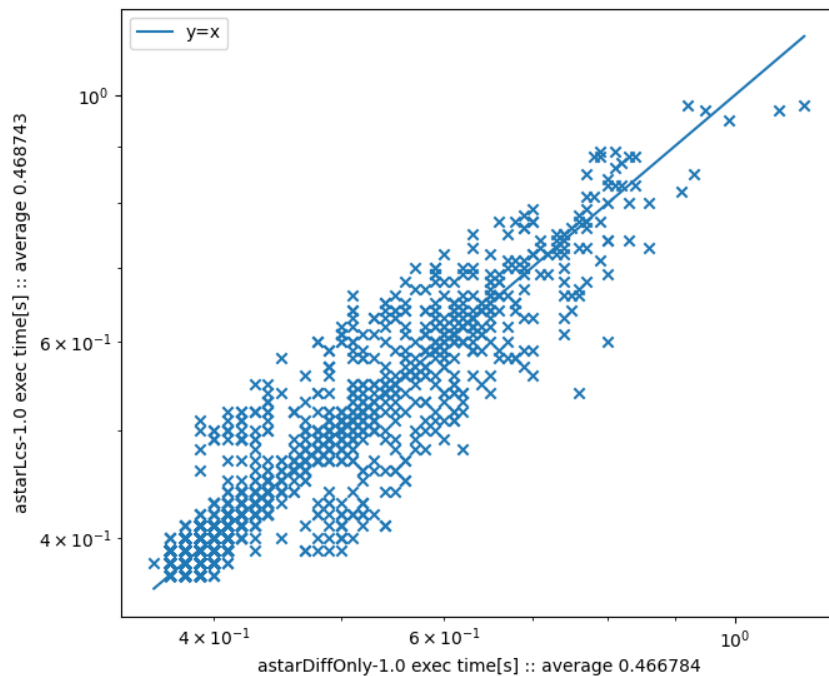
Reduced Searched Node :: 80%

Reduced Time :: 0%

commit=d0db4aececb5e6865b67f9dd2a46053e38cc4848
N = 1710



commit=d0db4aececb5e6865b67f9dd2a46053e38cc4848
N = 1710

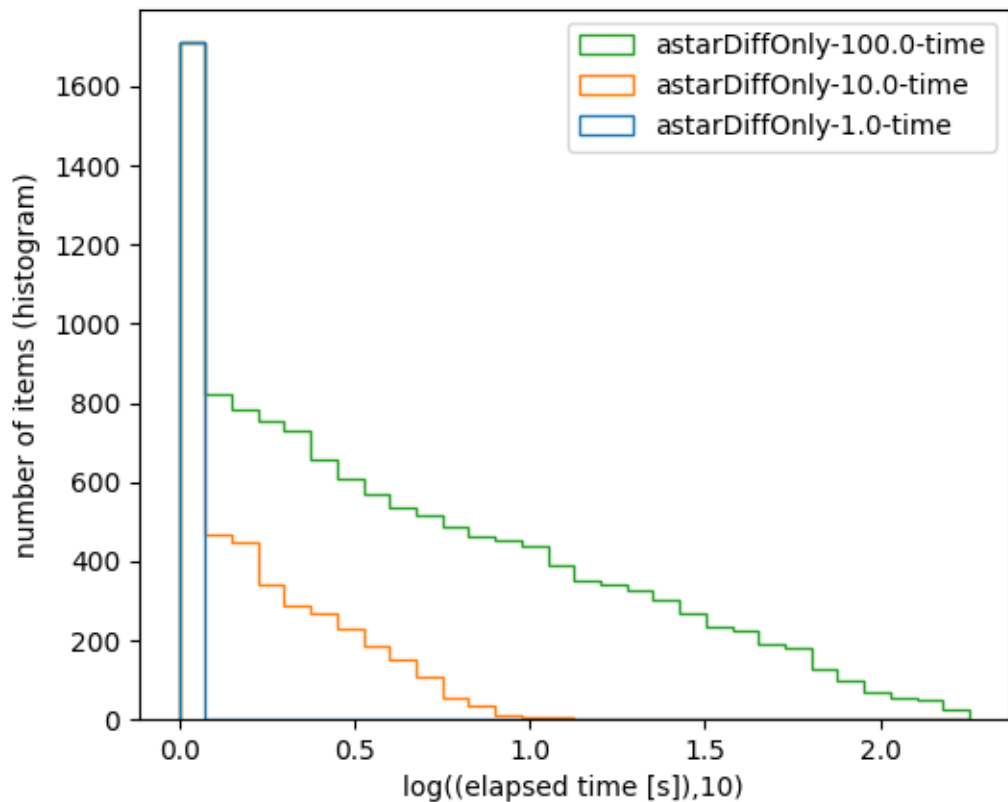


Appendix:

Best parameter for stack limit L

commit=d0db4aeceb5e6865b67f9dd2a46053e38cc4848

N = 1710



Tested on $L :=$ 10000.0,
100000.0,
1000000.0

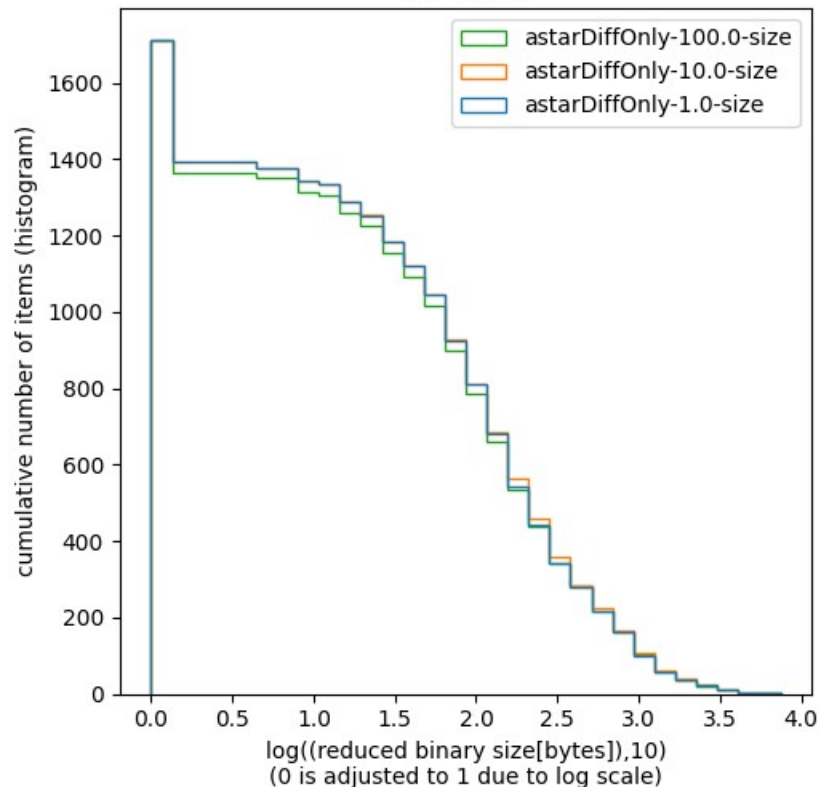
Execution time seems to
be proportional to L .

Appendix:

Best parameter for stack limit L

commit=d0db4aececb5e6865b67f9dd2a46053e38cc4848

N = 1710



Increasing L
scarcely reduces code size.

i.e.

Speeding up optimization is
not so urgent.

Appendix: Future work

- Extend target ops of exhaustive search
 - Data type constructing / deconstructing ops
PAIR / CAR / CDR / CONS
 - Calculation ops
ADD / SUB / MUL / DIV / LE
- Faster algorithm for exhaustive search